

No Idle Cores

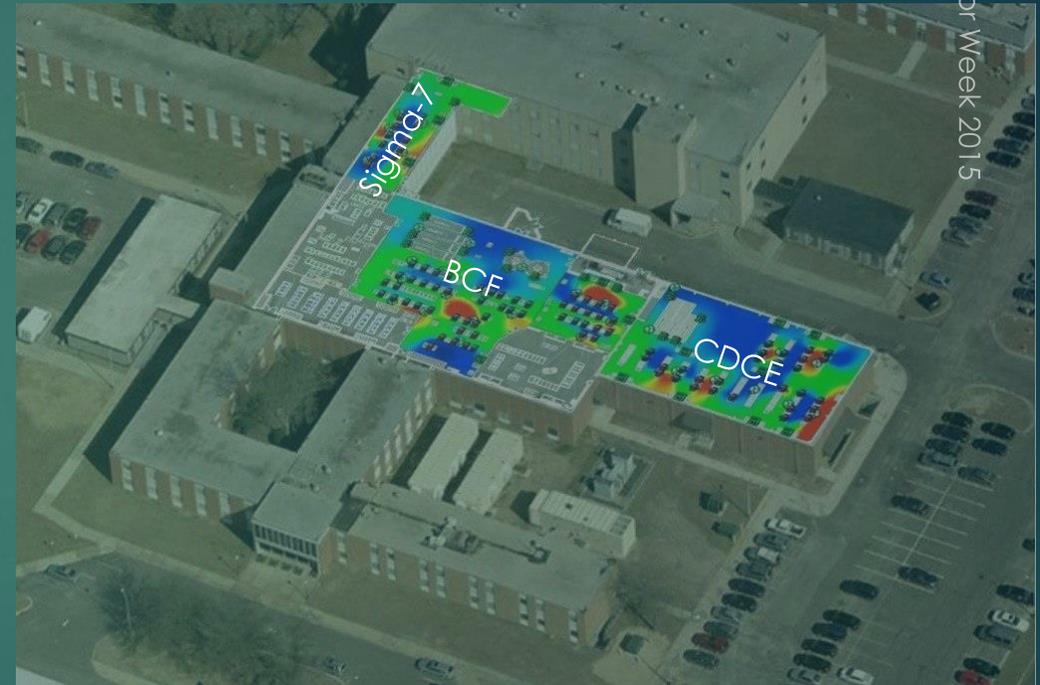
DYNAMIC LOAD BALANCING IN ATLAS & OTHER NEWS

WILLIAM STRECKER-KELLOGG <WILLSK@BNL.GOV>

RHIC/ATLAS Computing Facility Overview

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- ▶ Main HTCondor Pools
 - ▶ STAR, PHENIX, ATLAS
 - ▶ Each just over 15kCPU
- ▶ Running stable 8.2.7
- ▶ RHIC Pools
 - ▶ Individual+Special Users
 - ▶ Workload management done by experiments
- ▶ ATLAS: Focus of this talk
 - ▶ Local batch systems driven by external workload manager (PANDA)
 - ▶ Jobs are pilots
 - ▶ Scheduling→provisioning



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ATLAS Configuration

- ▶ Use Hierarchical Group Quotas + Partitionable Slots
 - ▶ My [HTCondor Week talk](#) last year was all about this
- ▶ A short recap:
 - ▶ PANDA Queues map to groups in a hierarchical tree
 - ▶ Leaf-nodes have jobs
 - ▶ [Surplus-sharing is selectively allowed](#)
- ▶ Group allocation controlled via web-interface to DB
 - ▶ Config file written when DB changes

- ▶ All farm has one STARTD config

```
SLOT_TYPE_1=100%  
NUM_SLOTS=1  
NUM_SLOTS_TYPE_1=1  
SLOT_TYPE_1_PARTITIONABLE=True  
SlotWeight=Cpus
```



Partitionable Slots

- ▶ Each batch node is configured to be partitioned into arbitrary slices of CPUs
 - ▶ Condor terminology:
 - ▶ **Partitionable slots** are automatically sliced into **dynamic slots**
- ▶ Multicore jobs are thus accommodated with no administrative effort
 - ▶ Only minimal (~1-2%) defragmentation necessary
 - ▶ Empirically based on our farm—factors include cores/node, job sizes & proportions, and runtimes. Something like

$$\text{draining} = (\text{job-length} * \text{job-size}^2) / (\text{machine-size} * \% \text{mcore} * \text{occupancy})$$

Defragmentation Policy

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Defragmentation Daemon

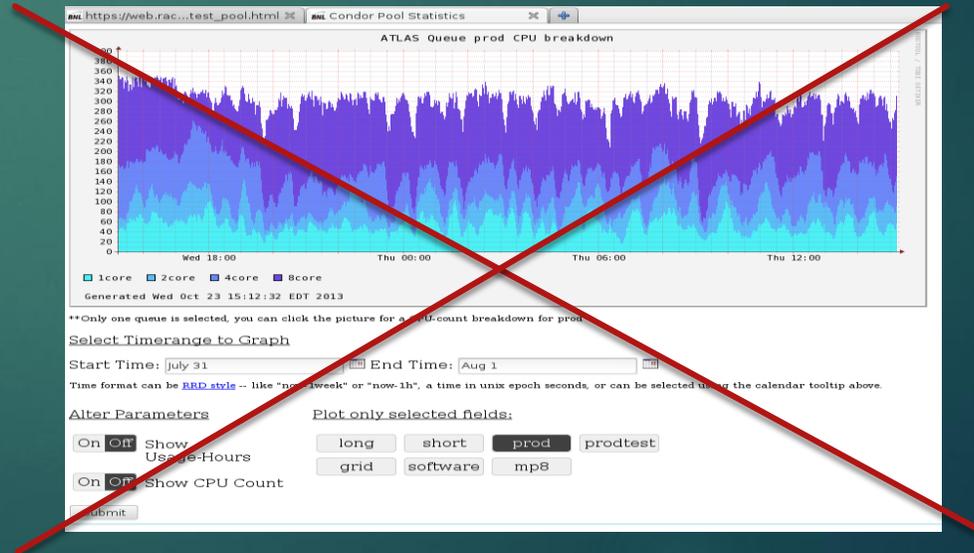
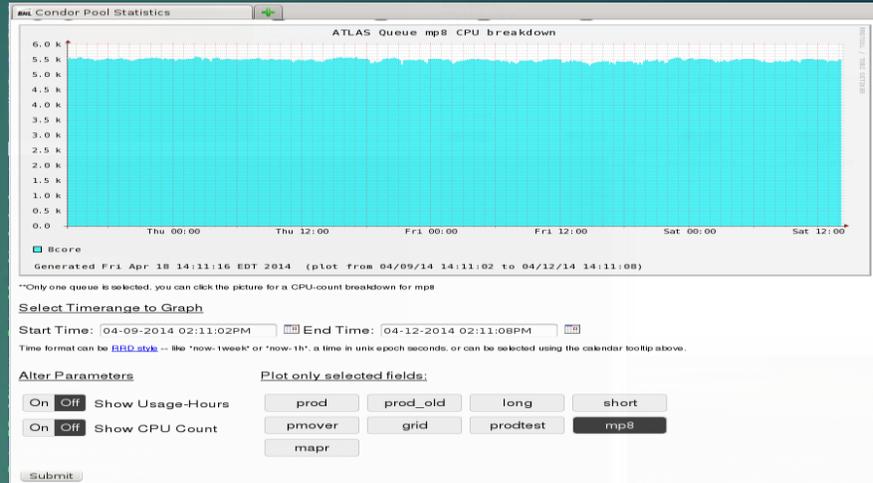
- ▶ Start Defragmentation
 - ▶ `(PartitionableSlot && !Offline && TotalCpus > 12)`
- ▶ End Defragmentation
 - ▶ `(Cpus >= 10)`
- ▶ Rate: max 4/hr

Key change: Negotiator Policy

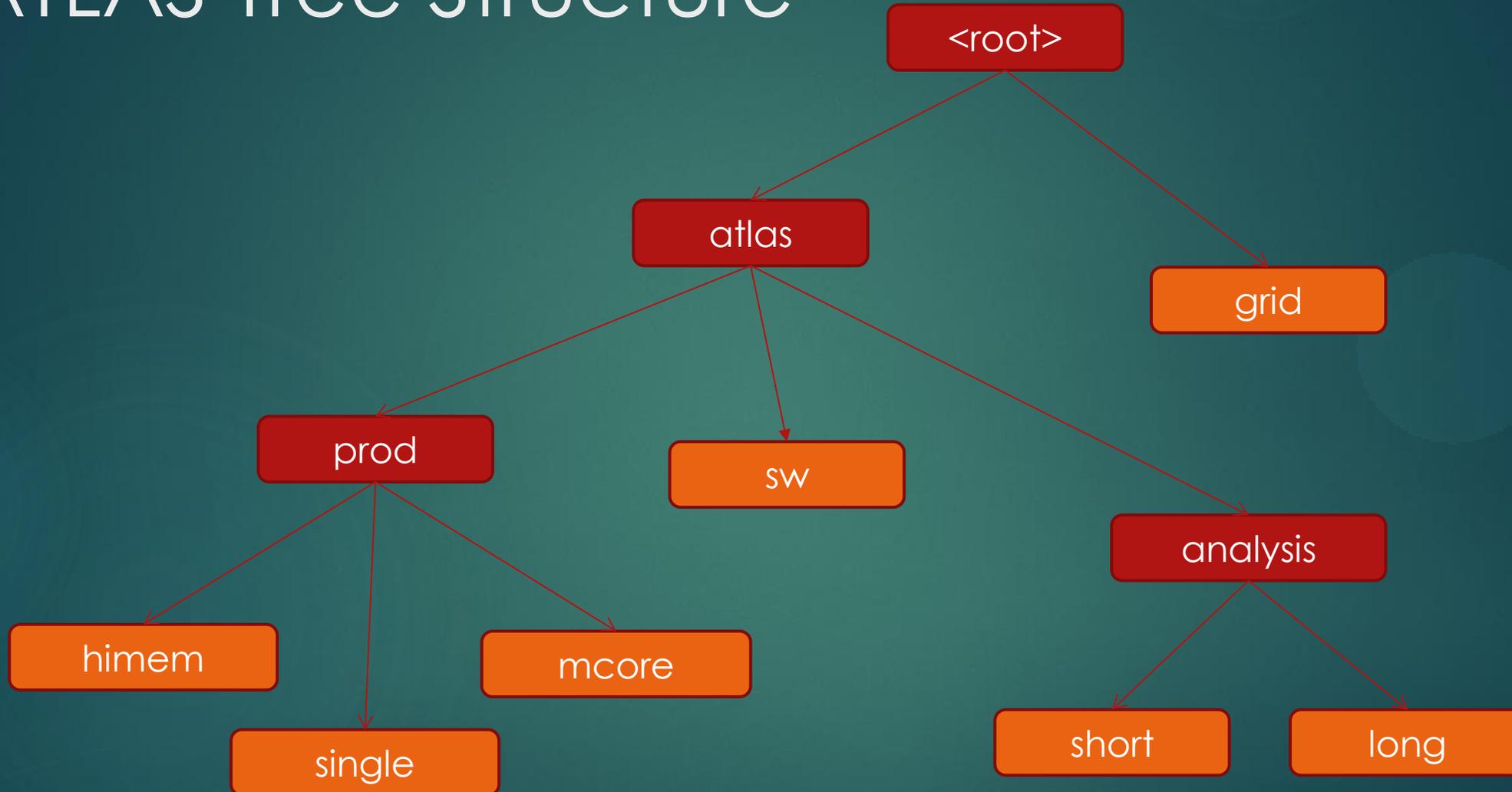
- ▶ Setting `NEGOTIATOR_POST_JOB_RANK`
- ▶ Default policy is breadth-first filling of **equivalent** machines
 - ▶ `(Kflops - SlotId)`
- ▶ Depth-first filling preserves continuous blocks longer
 - ▶ `(-Cpus)`

PANDA Queues

- ▶ PANDA Queues
 - ▶ One species of job per-queue
 - ▶ Map to groups in our tree
- ▶ Currently two non-single-core queues
 - ▶ 8-core ATHENA-MP
 - ▶ 2-core (Actually high-memory)
 - ▶ No support yet for SlotWeight!=cpus
 - ▶ Have 2Gb/core, so 4Gb jobs get 2 cores



ATLAS Tree Structure

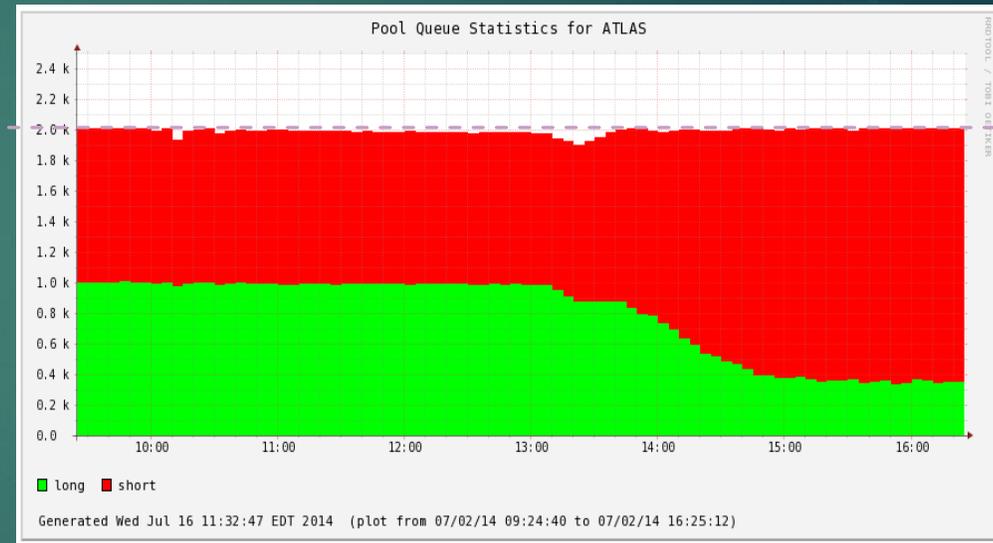


Surplus Sharing

- ▶ Surplus sharing is controlled by boolean *accept_surplus* flag on each queue
 - ▶ Quotas are normalized in *units of SlotWeight (CPUs)*
- ▶ Groups with flag set to True can take unused slots from their siblings
 - ▶ Parent groups with flag allow surplus to “flow down” the tree from their siblings to their children
 - ▶ Parent groups without *accept_surplus* flag constrain surplus-sharing to among their children

Surplus Sharing

- ▶ Scenario: **analysis** has quota of 2000 and no *accept_surplus*; **short** and **long** have a quota of 1000 each and *accept_surplus* on
 - ▶ short=1600, long=400...possible
 - ▶ short=1500, long=700...impossible (violates analysis quota)



Where's the problem? (it's starvation)

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- ▶ Everything works perfectly with all single-core, just set `acctpet_surplus` everywhere!
- ▶ However... Multicore jobs will not be able to compete for surplus resources fairly
 - ▶ Negotiation is greedy, if 7 slots are free, they won't match an 8-core job but will match 7 single-core jobs in the same cycle
 - ▶ If any multicore queues compete for surplus with single core queues, the multicore will always lose
- ▶ **A solution outside Condor is needed**
 - ▶ Ultimate goal is to maximize farm utilization—No Idle Cores!

Dynamic Allocation

- ▶ A program to look at the current state of the demand in various queues and set the surplus-flags appropriately
 - ▶ Based on comparing “weight” of queues
 - ▶ Weight defined as size of jobs in queue (# cores)
 - ▶ Able to cope with any combination of demands
 - ▶ Prevents starvation by allowing surplus into “heaviest” queues first
 - ▶ Avoids both single-core and multicore queues competing for the same resources
 - ▶ Same algorithm is extensible up the tree to allow sharing between entire subtrees
 - ▶ Much credit to Mark Jensen (summer student in '14)

Balancing Algorithm

- ▶ Groups have the following properties pertinent to the algorithm
 - ▶ Surplus flag
 - ▶ Weight
 - ▶ Threshold
 - ▶ Demand

- ▶ If Demand > Threshold a queue is considered for sharing

```
MariaDB [group_quotas]> select group_name, quota, weight, accept_surplus as surplus, surplus_threshold as threshold from atlas_group_quotas;
```

group_name	quota	weight	surplus	threshold
group_atlas	14823	2	0	0
group_atlas.analysis	4008	1	0	0
group_atlas.analysis.long	2000	1	0	250
group_atlas.analysis.mcore	8	0	0	20
group_atlas.analysis.short	2000	1	0	250
group_atlas.prod	10810	2	0	0
group_atlas.prod.mp	6110	8	1	40
group_atlas.prod.production	3700	1	0	160
group_atlas.prod.test	1000	2	0	80
group_atlas.software	5	0	0	1
group_gridods	500	1	0	50

```
11 rows in set (0.00 sec)
```

```
MariaDB [group_quotas]>
```

Balancing: Demand

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- ▶ PANDA Queues are monitored for “activated” jobs
 - ▶ Polling every 2 minutes
- ▶ Last hour is analyzed
 - ▶ Midpoint sampling
 - ▶ Moving average
 - ▶ Spikes smoothed out
 - ▶ Queue considered “loaded” if calculated demand $>$ threshold

Extending Weight & Demand

- ▶ Leaf groups' weights are the cores they need (8, 2, 1)
- ▶ How to extend beyond leaf-groups?
 1. Define in custom priority order
 2. Define as `sum()` or `avg()` of child groups's weights
 - ▶ Problem with 1. is you can't guarantee starvation-free
 - ▶ For now, manually set weights to match what would be the case for 2.
- ▶ For demand and threshold: easy—sum of child-groups values

Balancing: Weight

<weight>/<threshold>

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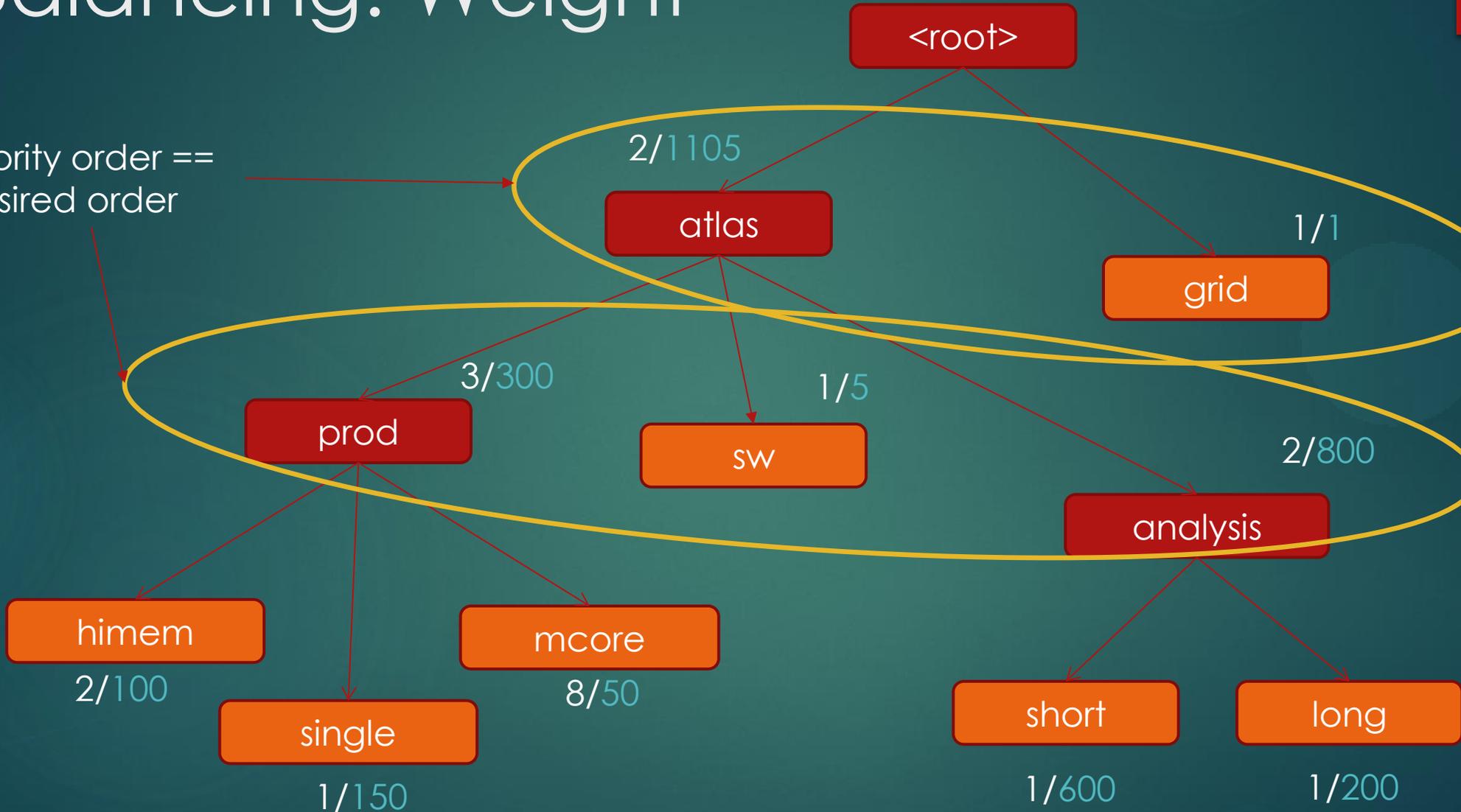
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Balancing: Weight

<weight>/<threshold>

Priority order ==
desired order



Algorithm

▶ The following algorithm is implemented

1. For each sibling-group in DFS order:
 1. For each member in descending weight order
 1. Set to TRUE unless it does **not** have demand **and** lower-weight groups **do**
 2. Break if set to TRUE

▶ In other words...

In each group of siblings, set *accept_surplus* to TRUE for all the highest-weighted groups that have demand

Balancing: All Full

Surplus ON

This is the same as having no surplus anywhere!



Balancing: No mc/hi

Surplus ON



Balancing: No prod

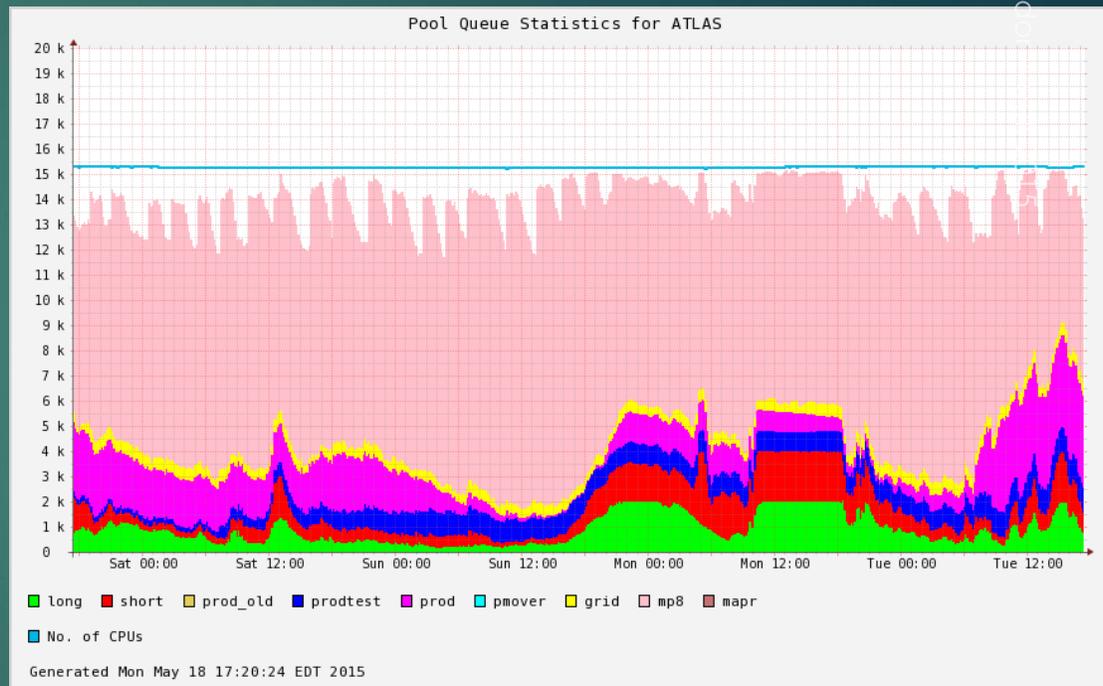
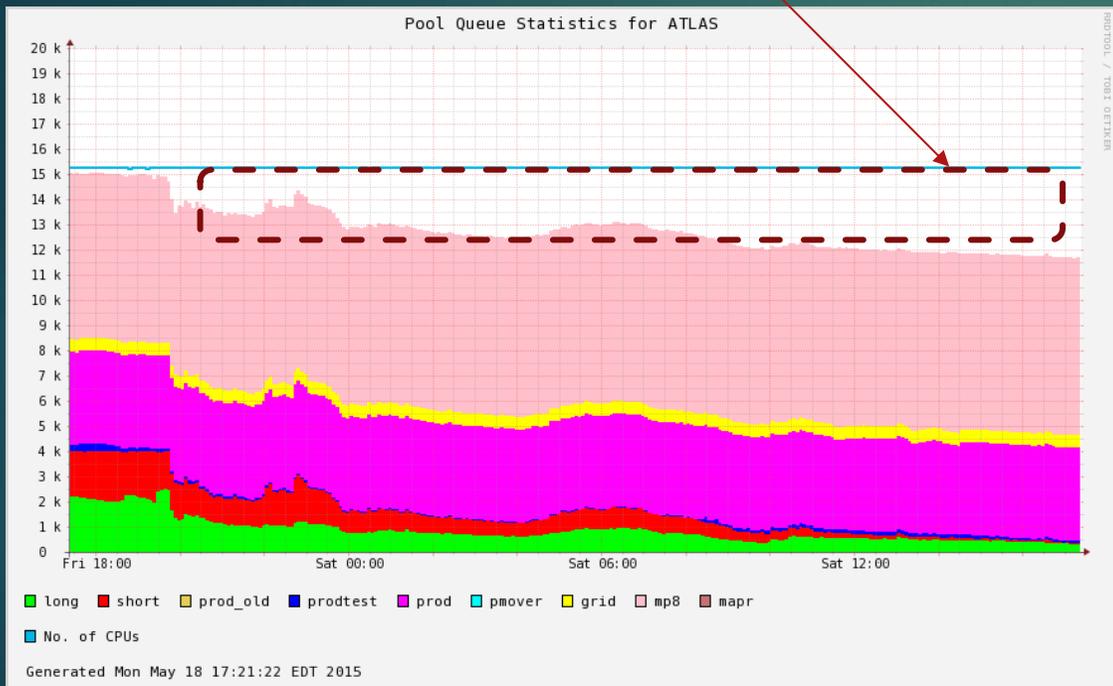
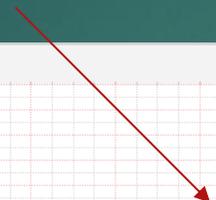
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Surplus ON



Results

Wasted Slots



Results & Context

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- ▶ Multicore is ready to take slack from other production queues
- ▶ Spotty analysis-demand the past few months has allowed many millions of CPU-hours to go unwasted
- ▶ If all ATLAS has a lull in demand, OSG jobs can fill the farm
 - ▶ Caveat: Preemption!
- ▶ Fast Negotiation
 - ▶ Averages for last 3 days:

Matches	14.99
Duration	7.05s

- ▶ Who is this useful for?
 - ▶ Algorithm works for any tree
 - ▶ Extensible beyond ATLAS where work is *structured* outside of batch system
 - ▶ A multi-tenant service provide with a hierarchy of priorities
 - ▶ Really a problem of efficient *provisioning*, not scheduling
- ▶ Constraints
 - ▶ Workflow defined outside of HTCondor
 - ▶ Segregation of multicore in separate queues for scheduling

Desired Features & Future Work

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Preemption

- ▶ Wish to maintain reasonably minimum-runtime to grid jobs
- ▶ When ATLAS demand comes back, need OSG jobs to be evicted
- ▶ Require preempting the *dynamic* slots that are created under the *partitionable* one
 - ▶ Work is progressing along these lines, although final state is not clear

SlotWeight != CPUs

- ▶ Would like to “value” RAM less than CPUs for jobs
 - ▶ High-memory kludge is inelegant
 - ▶ Not extensible to different shaped jobs (high-RAM/low-CPU, vice versa)
- ▶ Tricky because total slot-weight of the farm needs to be constant to give meaning to quota allocation

The End

QUESTIONS? COMMENTS?

THANKS TO MARK JENSEN, AND THE HTCONDOR TEAM!