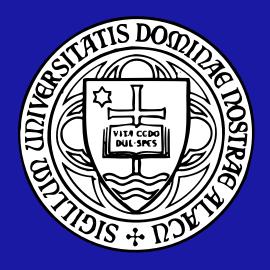
The LAM Implementation of MPI: Features, Dynamic Process Control, and Checkpointing

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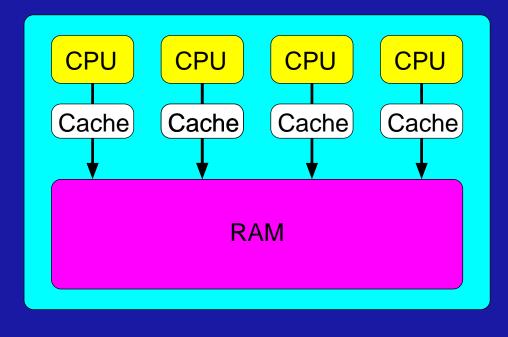


Overview

- Introduction: Parallel Computing
- MPI (and others)
- The LAM implementation of MPI
- LAM + Condor = Lamdor
- Conclusions / Future Work

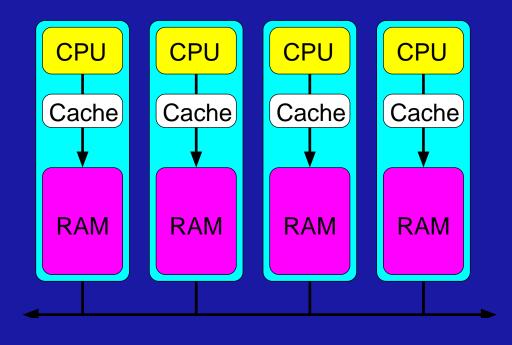
Introduction: Parallel Computing

- Shared memory
 - Typically multi-threaded, sometimes multi-process
 - All sharing common memory
 - Not [directly] the focus of this talk



Introduction: Parallel Computing

- Distributed memory
 - Typically separate processes
 - Explicit sending of messages; little [direct] use of shared memory



Message Passing: The Contenders

- Parallel Virtual Machine (PVM)
 - Research project at Oak Ridge National Labs
 - First message passing package on clusters
 - Attracted a *large* user base
- Message Passing Interface (MPI)
 - MPI-1 standardized in 1994, MPI-2 standardized in 1997
 - Vendor and open source implementations
 - Source code portable

PVM: The Good

- Daemon-based run-time environment
- Popularized manager / worker model using dynamic processes
- Load factoring / environment querying
- Inter-implementation communication
- Still has a large user following

PVM: The Bad

- Usually forces an extra buffer copy
- No true asynchronous communication
- Nondeterministic behavior (particularly with groups)
- Weak message safety (only one context at a time)
- Losing vendor support

What is MPI?

- A specification for a message passing API
- Two documents:
 - MPI-1: Basic message passing (send, receive, collectives, etc.)
 - MPI-2: Extensions to MPI-1 (one-sided, C++, dynamic processes, etc.)
- Specifically written to enable high performance
- Designed for clusters all the way up to "Big Iron"

MPI: The Good

- Learned from previous designs: NX, Zipcode, PVM, etc.
- No extra memory copy
- Capable of true asynchronous communication
- Deterministic behavior, ease of discovering identity
- Strong message safety
- Strong (and continuing) vendor support
- Uses fastest message passing available (shmem, TCP, etc.)

MPI: The Bad

- Some of the previous is "implementation dependent"
 - True asynchronous communication
 - Use of fastest message passing channel
- No [portable] fault tolerance
- MPI's design does not preclude any of these, but much of this is [intentionally] left unspecified

MPI Terminology

- Rank: A single entity in a parallel job
- Communicator: A group of ranks plus a unique message passing context
- MPI_COMM_WORLD: Default communicator that contains all ranks

MPI Code Example: Hello World

```
int main(int argc, char* argv[]) {
int me, total;
MPI_Init(&argc, &argv);
MPI_Comm_rank(MPI_COMM_WORLD, &me);
MPI_Comm_size(MPI_COMM_WORLD, &total);
```

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The LAM Implementation of MPI

- Originally written at the Ohio Supercomputing Center
 - Targeted at transputers
 - Top MPI layer was added later
 - MPI has since become the main focus of work
- Everyone graduated, moved on
 - LAM/MPI was orphaned for about a year
- Project moved to Notre Dame in 1998

MPI Conformance

- Full MPI-1 conformance
- Much MPI-2 functionality
 - MPI I/O (ROMIO)
 - MPI C++ bindings
 - One-sided communication
 - Dynamic process control
- Interoperable MPI (IMPI)
 - Point-to-point and some collectives

LAM/MPI: Features

- Daemon-based run-time environment
 - Fast startup of user programs
 - Guaranteed cleanup of user programs
 - External monitoring
- Flexible mpirun
 - SMP-friendly syntax
 - SPMD or MPMD
 - Can distribute executables; no global filesystem required
 - Pseudo-tty support
 - Environment variable and working directory export

LAM/MPI: More Features

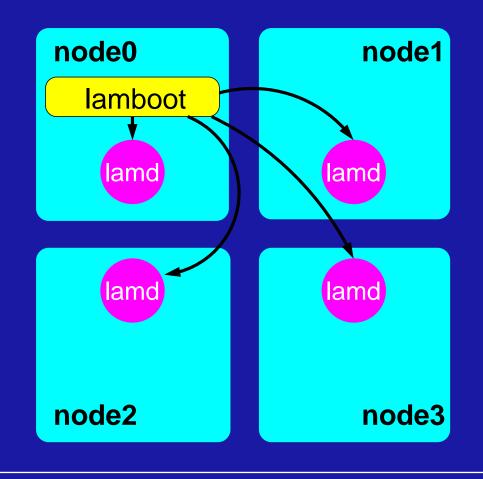
- Optimized point-to-point message passing
 - "Short-circuit" optimization
 - Short call stack; "uncomplicated" engine
 - Combined TCP / shared memory message passing
 - True asynchronous message passing
- Monitoring tools
 - XMPI: GUI message passing patterns
 - mpimsg: pending messages
 - mpitask: running LAM tasks

LAM/MPI: Still More Features

- Heterogeneous support
 - Portable to most POSIX systems
 - On-the-fly endian conversion (if necessary)
- Debugging support
 - Ability to mpirun non-MPI executables (e.g., debuggers)
 - Totalview support on the way
 - Purify clean
 - Open source (some users actually *do* source dive!)
 - Lots of online help: web pages, man pages

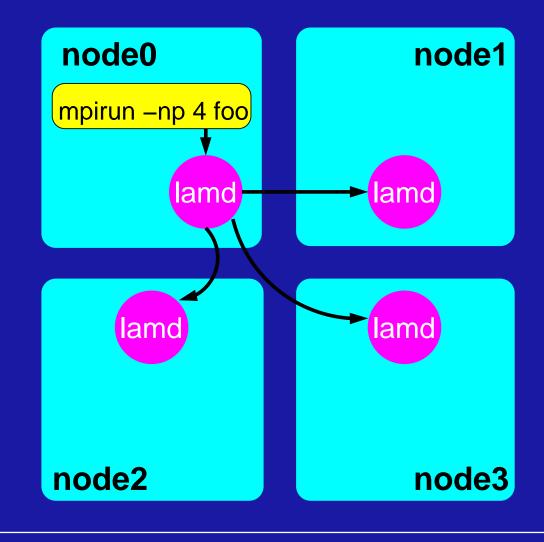
Daemon-Based Run-Time Environment

- User level, not root level
- Launch the LAM RTE: lamboot <hostfile>



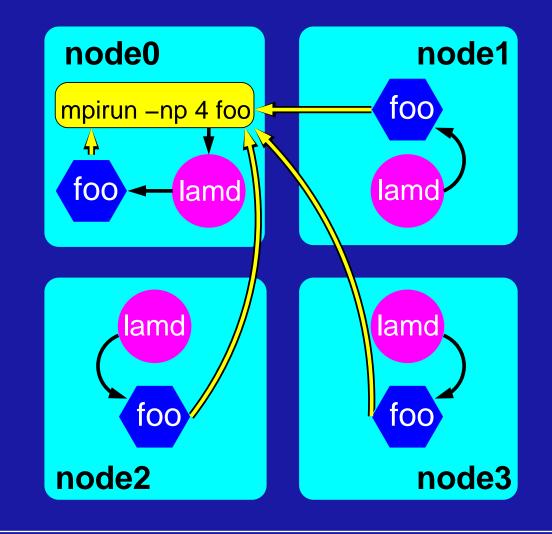
Running MPI Programs

• mpirun sends a message to the local daemon



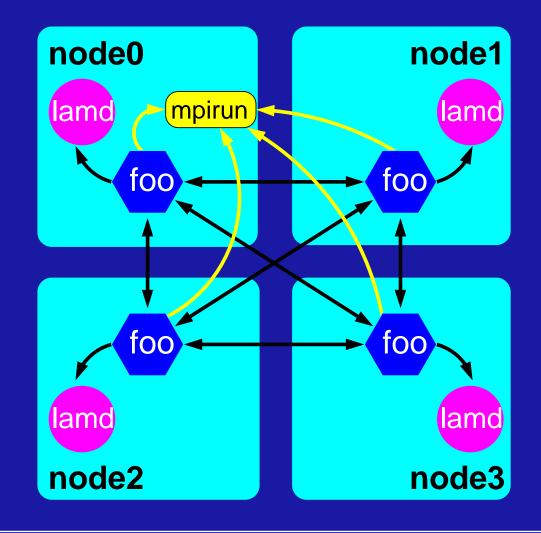
Running MPI Programs: Step 2

• The daemons fork / exec the child, setup stdin / stdout, etc.



Running MPI Programs: Step 3

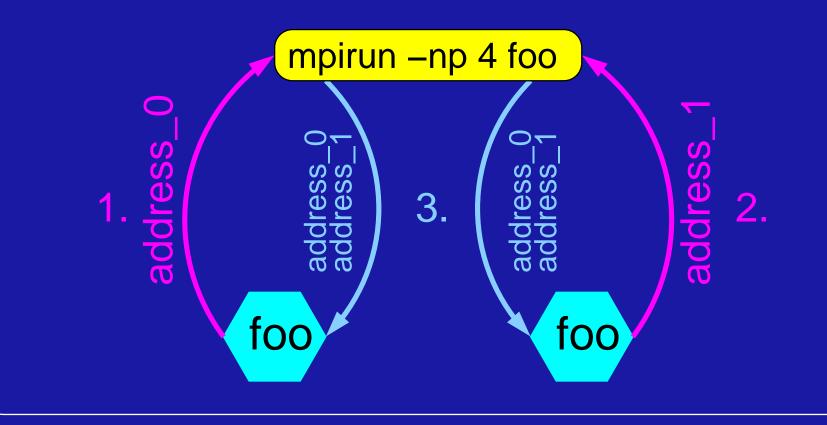
• During MPI_Init(), each MPI rank connects all others



MPI_Init: Mutual Awareness

Out-of-band messaging is used during MPI_Init()

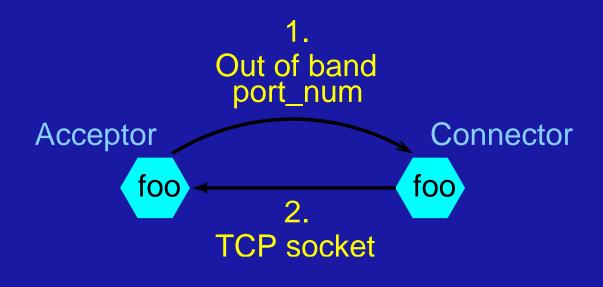
- Each rank contacts mpirun
- mpirun sends full list of out-of-band peer addresses



MPI_Init: Peer-to-Peer Setup

• More out-of-band messaging used between MPI ranks

- Each rank opens a dynamic "listening" socket
- Pairwise, "acceptor" rank sends port number to "listener"

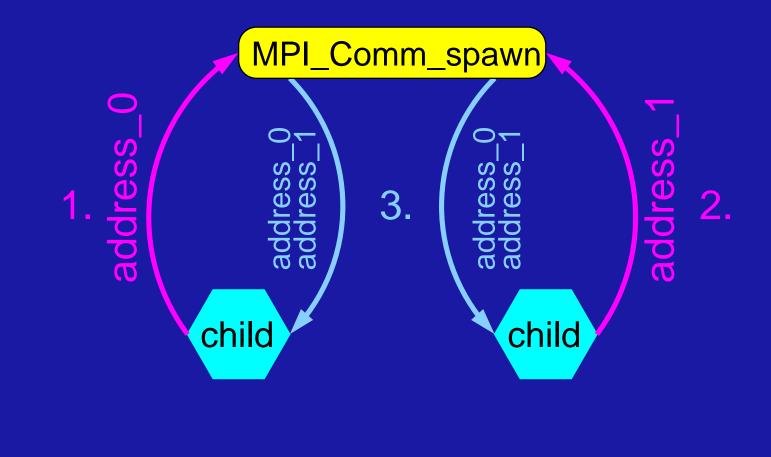


MPI Dynamic Processes

- MPI_Comm_spawn() is used to launch a group of children processes
 - It is a collective (blocking) call across the spawning communicator
- Children processes will have their own unique MPI_COMM_WORLD
- Parents and children will share a "bridge" communicator that they can communicate with

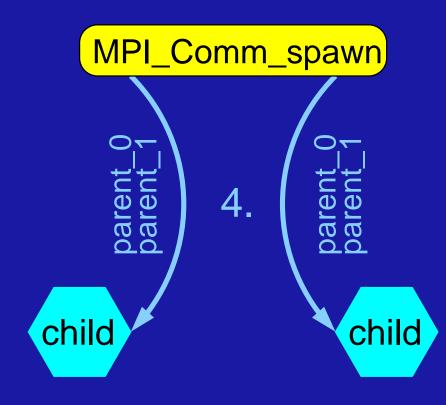
MPI_Comm_spawn: Same as MPI_Init

- Replace mpirun instance with MPI_Comm_spawn() instance
- Do same out-of-band messaging



MPI_Comm_spawn: Slight Differences

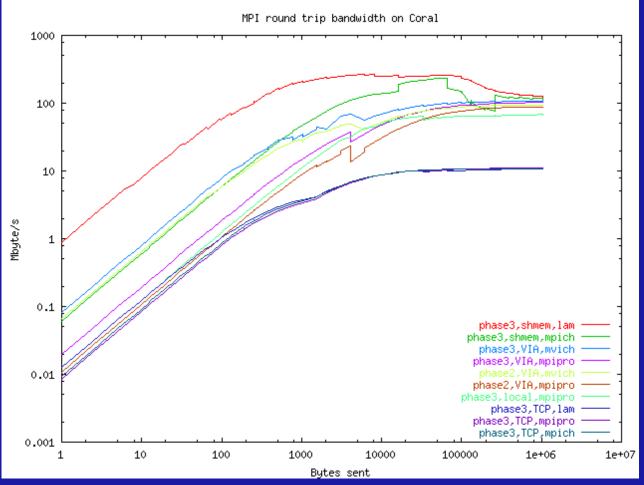
- Children must create parent communicator
- MPI_Comm_spawn() must also send parent addresses



MPI Bandwidth Performance: TCP

- MPI performance on the Coral cluster
- ICASE / NASA Langley Research
- 32 nodes
 - Phase 2: 16 dual pentium III/500, PC100
 - Phase 3: 16 dual pentium III/800, PC133
- Gigabit ethernet interconnection network

MPI Bandwidth Performance: TCP

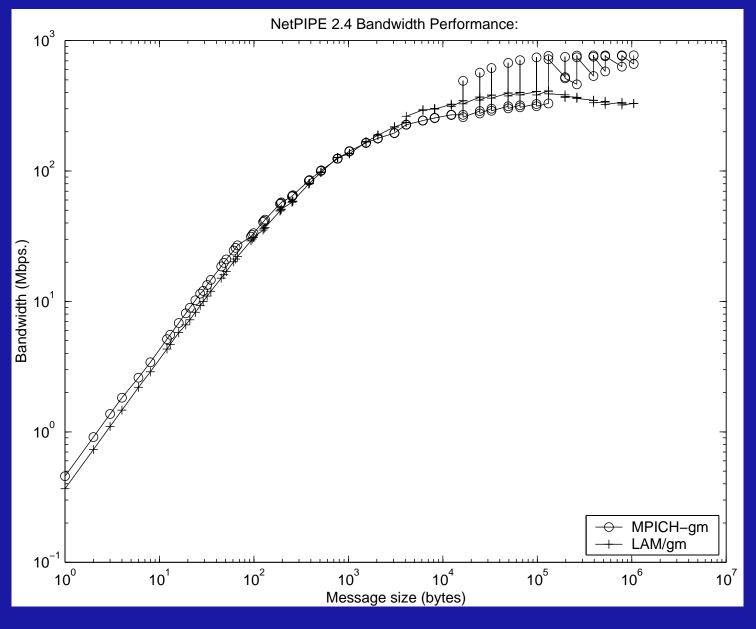


- MPICH latency over TCP: 175us
- LAM latency over TCP: 129us

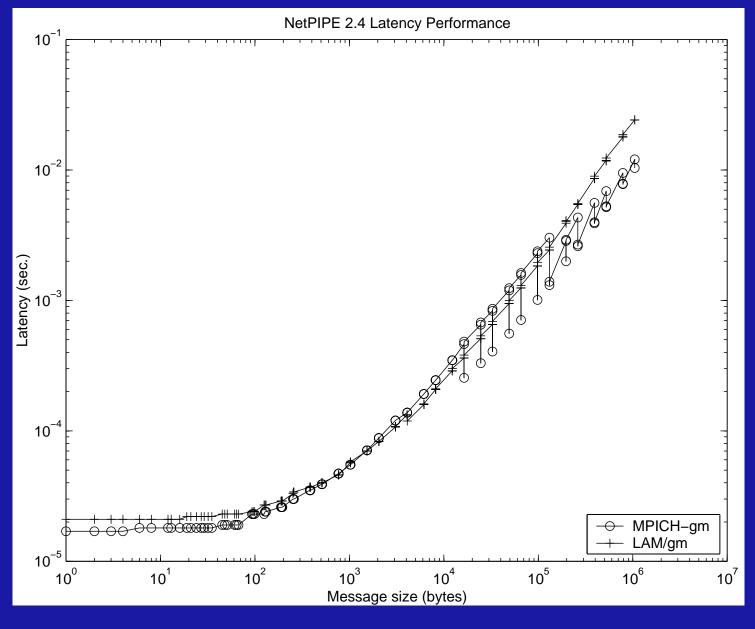
MPI Bandwidth Performance: Myrinet

- Linux pentium cluster at Indiana University
- Myrinet interconnection network
- Numbers collected using NetPIPE 2.4 benchmark application

MPI Bandwidth Performance: Myrinet



MPI Latency Performance: Myrinet



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LAM + Condor = Lamdor

• We want to run LAM jobs under Condor

- Get MPI-2 features in Condor (spawn, put / get, etc.)
- Take MPI support burden away from Condor team
- Bring MPI into distributed [dynamic] computing
 - Manager / worker process model
 - Some degree of fault tolerance
- Eventual goal: cycle-scavenging parallel jobs
 - Transparent checkpoint/migrate/restart support for MPI jobs

Applies Outside of Condor Environments

- Clusters aren't as stable as Beowulfers would have you believe
 - Nodes die, switches and routers fail, etc.
 - That is: even dedicated clusters are dynamic clusters
- Scale LAM up to grid-sized problems: dynamic MPI environments
- Use the standalone checkpoint library outside of a Condor flock
 - Resource maintenance (e.g., Beowulf-style clusters)
 - Queue management / dedicated resource time
 - Program fault protection

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Conclusions / Future Work

- Get LAM/MPI to run under the Condor static scheduler
- Extend LAM/MPI to dyanmic and unreliable environments
 - Properly define (redefine?) MPI semantics / models for dynamic environments
- Extend LAM/MPI to handle grid-sized problems
- Experiment with checkpoint / migrate / restart schemes