WALi User Manual

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1 Introduction

WALi¹ is an open source library implementation of a Weighted Pushdown System (WPDS). WPDSs have been shown to be a powerful formalism for performing interprocedural dataflow analysis [3, 1, 2]. (For a listing of relevant papers see http://www.cs.wisc.edu/wpis/wpds.)

2 Installation

WALi has been compiled on Linux, Cygwin+Windows/XP and using Visual Studio 2005. In either case, one must first download the WALi source tree. The latest can be found http://www.cs.wisc.edu/wpis/wpds/wali/WALi-latest.tar.gz. This link will download a zipped tar file of the WALi source tree.

2.1 Linux and Cygwin

WALi requires a working Python installation², the SCons build tool, and the common utility program curl. The following commands will download and extract the WALi source tree.

```
$ cd <place to untar WALi>
$ curl -0 http://www.cs.wisc.edu/wpis/wpds/wali/WALi-latest.tar.gz
$ tar zxvf WALi-latest.tar.gz
```

We will use the name WALiDir to denote the top level directory of the WALi source tree. It is convention to define an environment variable \$WALiDir that holds this value.

```
$ cd WALiDir
$ scons
```

 $^{^1 \}rm WALi$ is short for "Weighted Automaton Library" $^2 \rm Only$ Python v2.5 has been fully tested

This will compile and link the WALi library in the directory WALi-1.2/lib. (For cygwin, a static library is built.)

The WALi source tree comes bundled with some examples and addons. These can be compiled with the commands scons examples and scons addons, respectively. The examples include weight domains (§4) for reachability, kill-gen problems, and affine-relations. The addons includes C++ code for parsing WPDS queries specified in XML format. The parsing code makes use of the Apache xerces-c XML parsing libarary.

After compiling the library, the directory \$WALiDir/lib should be added to the environment variable LD_LIBRARY_PATH or PATH for Linux or Cygwin, respectively.

2.2 Visual Studio

Download and extract the WALi source tree. Underneath the top-level directory of the WALi source tree, there is a directory named Projects. Underneath the Projects directory there is a file WALi.vs80.sln, which is a Visual Studio solution file containing three projects: WALi, LiveVar and Parse. To incorporate WALi into an existing Visual Studio project, import the WALi project (WALi.vcproj) if possible, or copy the settings over. Else, compile the WALi library using the WALi solution and set the include and library directories for an existing Visual Studio project to the appropriate place.

3 Interprocedural Control-Flow Graph Encoding

For performing dataflow analysis, the standard practice is to encode the program's interprocedural control flow graph as a single state PDS (see Fig. 1). The weights that annotate the rules of the PDS are dataflow transformers that encode the effect of (abstractly) executing a program statement associated with the rule.

Rule	Control flow modeled
$(p, n_1) \hookrightarrow \langle p, n_2 \rangle$	Intraprocedural edge $n_1 \rightarrow n_2$
$\langle p, n_c \rangle \hookrightarrow \langle p, e_f \ r_c \rangle$	Call to f , with entry e_f , from n_c that returns to r_c
$\langle p, x_f \rangle \hookrightarrow \langle p, \epsilon \rangle$	Return from f at exit x_f

Figure 1: The encoding of an ICFG's edges as PDS rules.

4 Implementing a Weight Domain

A WALi user defines a weight domain *D* that encodes the desired abstract domain. *D* must be a subclass of the provided abstract class wali::SemElem. Inside of WALi, all instances of SemElem are reference counted. The reference counting implementation is defined by the C++ template class wali::ref_ptr<T>. For easier notation, WALi provides the type definition typedef wali::sem_elem_t wali::ref_ptr<wali::SemElem>. We next describe each of the methods that must be overridden by the class *D*.

```
4.1 One - 1
```

```
sem_elem_t one() const;
```

one returns an instance of the $\overline{1}$ element.

4.2 Zero - 0

sem_elem_t zero() const;

zero returns an instance of the $\overline{0}$ element.

4.3 Combine - \oplus

sem_elem_t combine(SemElem * se);

combine returns a new weight that is the combination of this and the parameter se.

4.4 Extend - \otimes

sem_elem_t extend(SemElem * se);

extend returns a new weight that is equal to this extended by the parameter se (this \otimes se). extend is typically related to functional composition. Is this regard, this \otimes t is functionally equivalent to t \circ this.

4.5 Equal

bool equal(SemElem * se) const;

equal returns true if two weights are equal and false if not. There is currently no method specifically designed to deal with partial orders. However, for any two semiring elements α and β , $\alpha \subseteq \beta \Leftrightarrow \alpha = (\alpha \oplus \beta)$.

4.6 Print

std::ostream & print(std::ostream & o) const;

print writes a semiring element to the passed in std::ostream parameter. It should return the same std::ostream when finished.

5 Examples

5.1 Reachability Weight Domain

The following weight domain implements simple reachability. The weight is $\overline{1}$ if it is reachable by the WPDS and $\overline{0}$ otherwise. The C++ header and source files are distributed with WALi under the Examples directory.

```
Code Listing 5.1 (Weight domain implementing reachability.).
#include "wali/SemElem.hpp"
```

```
using wali::SemElem;
using wali::sem_elem_t;
class Reach : public wali::SemElem
{
    public:
        Reach( bool b ) : isreached(b) {}
        virtual ~Reach() {}
        sem_elem_t one() const { return new Reach(true); }
        sem_elem_t zero() const { return new Reach(false); }
        // zero is the annihilator for extend
        sem_elem_t extend( SemElem* rhs ) {
          Reach* r = static_cast<Reach*>(rhs);
          return new Reach( isreached && r->isreached );
        }
        // zero is neutral for combine
        sem_elem_t combine( SemElem* rhs ) {
          Reach* r = static_cast<Reach*>(rhs);
          return new Reach( isreached || r->isreached );
        }
        bool equal( SemElem* rhs ) const {
          Reach* r = static_cast<Reach*>(rhs);
```

```
return isreached == r->isreached;
}
std::ostream & print( std::ostream & o ) const {
   return (isreached) ? o << "ONE" : o << "ZERO";
}
protected:
   bool isreached;
};</pre>
```

Using this weight domain is equivalent to using a Pushdown System without weights. All user created weights are $\overline{1}$ and unreachable configurations (abstractly) have weight $\overline{0}$.

6 Creating a Weighted Pushdown System

In this section, we show how to translate pseudo code following pseudo code is translated into a WPDS using the **Reach** semiring.

```
Code Listing 6.1 (Pseudo Code.).
```

```
// Pseudo Code //
x = 0
y = 0
fun f()
    nO: <$ f enter node $>
    n1: if(x = 0)
    n2:
            then y := 1
            else y := 2
    n3:
    n4: g()
    n5: <$ f exit node $>
fun g()
    n6: <$ g enter node $>
    n7: y := 4
    n8: x := 60
    n9: <$ g exit node $>
```

```
Code Listing 6.2 (WALi header files).
#include "wali/Common.hpp"
#include "wali/wpds/WPDS.hpp"
#include "wali/wfa/WFA.hpp"
#include "Reach.hpp"
```

First, a WPDS myWpds is created.

Code Listing 6.3 (Define the WPDS object myWpds.).
sem_elem_t reachOne(new Reach(true));
wali::wpds::WPDS myWpds;

Then the "keys" for the program locations are defined.

```
Code Listing 6.4 (Create Keys for program nodes.).
wali::Key p = wali::getKey("p");
wali::Key accept = wali::getKey("accept");
wali::Key n[10];
for( int i=0 ; i < 10 ; i++ ) {
    std::stringstream ss;
    ss << "n" << i;
    n[i] = wali::getKey( ss.str() );
}</pre>
```

The state and stack symbols of a WPDS rule have a type wali::Key. A key is a way of identifying a state or stack symbol of the WPDS. Each key has a unique wali::KeySource object associated with it. Some common sources have been defined like wali::StringSource and wali::IntSource. User's can define their own key source by subclassing the wali::KeySource class. The function wali::getKey is simply a helpful wrapper for creating a keys from C++ types std::string and int.

Once all the keys have been defined, the rules are added to the myWpds object.

```
Code Listing 6.5 (Add intraprocedural edges for f and g.). // f intraprocedural
```

```
myWpds.add_rule( p, n[0], p, n[1], reachOne);
myWpds.add_rule( p, n[1], p, n[2], reachOne);
myWpds.add_rule( p, n[1], p, n[3], reachOne);
myWpds.add_rule( p, n[2], p, n[4], reachOne);
myWpds.add_rule( p, n[3], p, n[4], reachOne);
myWpds.add_rule( p, n[6], p, n[7], reachOne);
myWpds.add_rule( p, n[6], p, n[7], reachOne);
myWpds.add_rule( p, n[8], p, n[9], reachOne);
myWpds.add_rule( p, n[8], p, n[9], reachOne);
myWpds.add_rule( p, n[4], p, n[6], n[5], reachOne);
// f calls g
myWpds.add_rule( p, n[5], p, reachOne);
```

// g return
myWpds.add_rule(p, n[9] , p , reachOne);

Then the initialized WPDS is printed to the standard output channel and marshalled as XML.

```
Code Listing 6.7 (Generic output methods.).
// Print the WPDS
myWpds.print( std::cout ) << std::endl;
// Marhasll the WPDS as an XML file
std::ofstream fxml( "myWpds.xml" );
myWpds.marshall( fxml );
fxml.close();</pre>
```

7 Querying the WPDS

WALi allows for two types of queries, *prestar* and *poststar*. A query takes as input a WPDS and a weighted finite automaton (WFA). A query outputs a new annotated WFA. WFAs are represented by the class wali::wfa::WFA. Transitions are added to the WFA class using the wali::wfa::WFA::addTrans method. The following sample code computes a *prestar* and *poststar* reachability query for the pseudo code (assuming the same objects are created as in the above C++ program).

```
Code Listing 7.1 (Prestar query.).
    wali::wfa::WFA prequery;
    prequery.addTrans( p, n[4], accept, reachOne );
    query.add_initial_state( p );
    query.add_final_state( accept );
    wali::wfa::WFA answer = myWpds.prestar(prequery);
    answer.print( std::cout );
```

```
Code Listing 7.2 (Poststar query.).
    wali::wfa::WFA postquery;
    postquery.addTrans( p, n[0], accept, reachOne );
    query.add_initial_state( p );
    query.add_final_state( accept );
    wali::wfa::WFA answer;
    myWpds.poststar(query,answer);
    answer.print( std::cout );
```

References

 T. Reps, S. Schwoon, S. Jha, and D. Melski. Weighted pushdown systems and their application to interprocedural dataflow analysis. *Science of Computer Programming*, 2005.

- [2] Thomas Reps, Akash Lal, and Nicholas Kidd. Program analysis using weighted pushdown systems. *FSTTCS*, 2007. Invited Paper.
- [3] Thomas W. Reps, Stefan Schwoon, and Somesh Jha. Weighted pushdown systems and their application to interprocedural dataflow analysis. In SAS, 2003.