

Preparing Object Code for Static Analysis

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Need for preprocessing

Data Dependence Analysis

```
int main(){
    int i,j, a[10];
    j=0;
    for(i=0;i<10;++i){
        a[i]=i,
    }
    return j;
}
```

A blue arrow points from the assignment `a[i]=i` to the question `Affects?`. A yellow arrow points from the question `No!` back to the assignment.

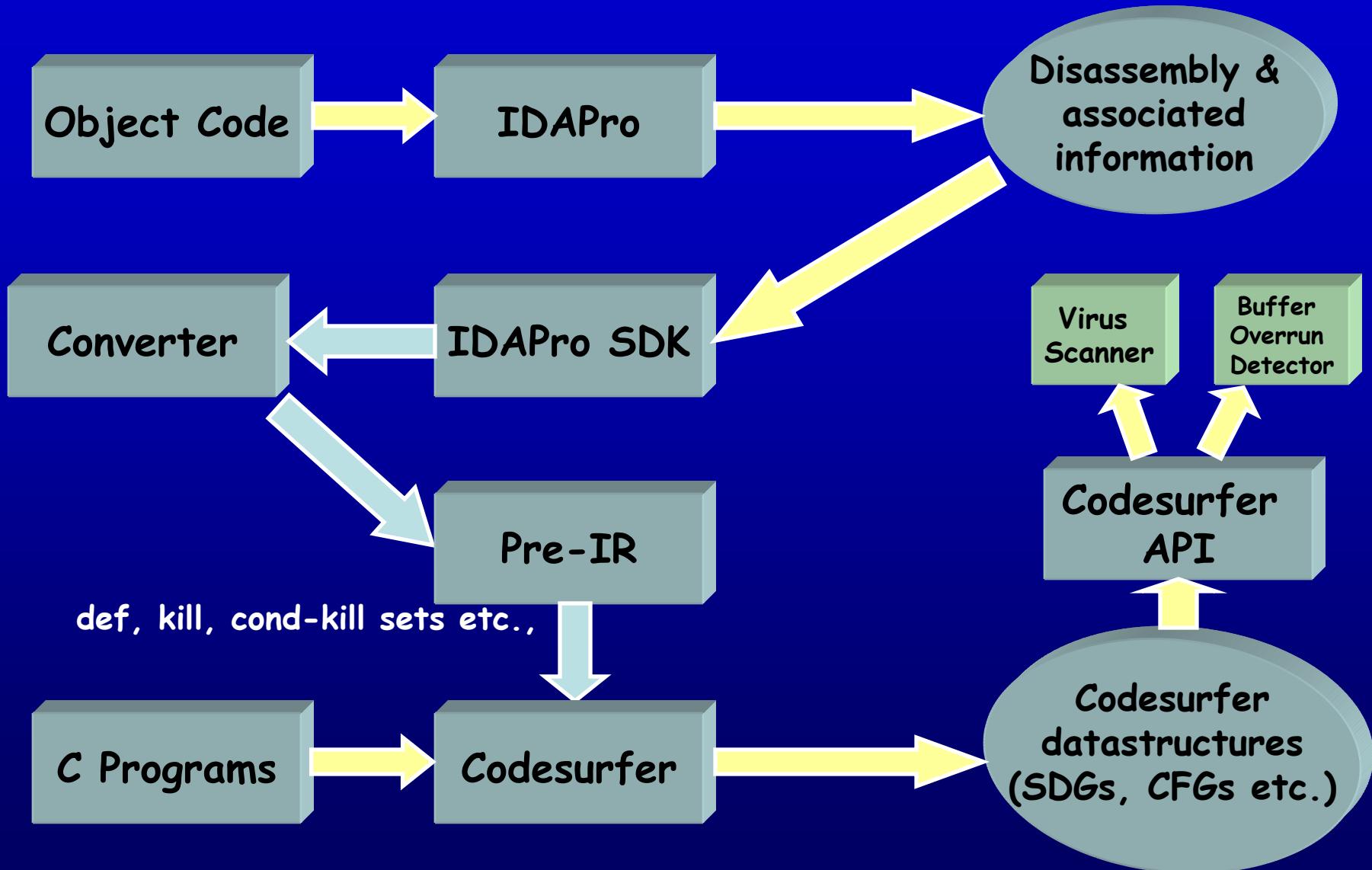
```
; ebx corresponds to variable i
sub    esp, 44
mov    [esp+40],0      ; j = 0
xor    ebx, ebx        ; i = 0
lea    ecx, [esp]
loc_9:
    mov    [ecx], ebx ; a[i]=i
    inc    ebx          ; i++
    add    ecx, 4
    cmp    ebx, 10       ; i<10?
    jl    short loc_9 ;
    mov    eax, [esp+40] ; return j
    add    esp, 44
    retn
```

A white arrow points from the question `Affects??` back to the assignment `eax, [esp+40]`.

Need for preprocessing

- In high level language programs
 - Variables
 - An abstract entity for memory
 - The entities on which the algorithm operates
 - We have a finite domain to operate on
- In object code
 - No properly defined entities
 - Has to be inferred

Existing infrastructure



Our Goal

- Discover the entities
- Annotate each program statement with def, kill and conditionally kill sets
- Feed it to Codesurfer
 - Already has a lot of static analysis algorithms implemented - Slicing, Chopping etc.,
- Can benefit
 - Virus scanner (Mihai)
 - Buffer Overrun Detector (Vinod)

Memory Model

Activation Records



- Four areas
 - Activation record
 - Global data
 - Heap
 - Expression stack
- Assumed to be disjoint
- Assumption should be validated

What are the entities?

Activation
Records



- Each area an entity?
 - Too inaccurate

What are the entities?

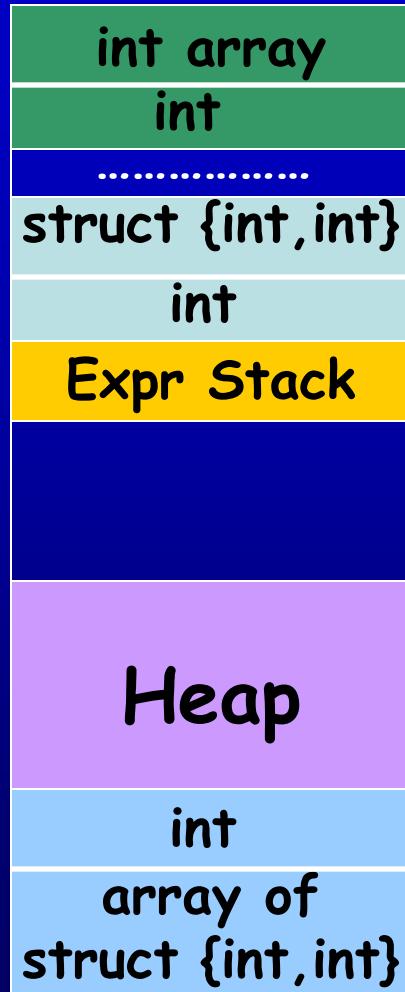
Activation
Records



- Each area an entity?
 - Too inaccurate
- Each byte an entity?
 - Accurate
 - But analysis is slow
 - 2^{32} addresses or more

What are the entities?

Activation Records



- Each area an entity?
 - Too inaccurate
- Each byte an entity?
 - Accurate
 - But analysis is slow
 - 2^{32} addresses or more
- Suppose we know the layout
 - Use the constituents as entities
 - Balanced solution

How to identify the entities?

- Aggregate Structure Identification
 - G. Ramalingam et al
- Algorithm
 - Ignores declarative information about aggregates (arrays, structs)
 - Decomposes aggregates - based on access patterns in program
 - Identified components - atoms
 - Unifies atoms which ought to have same type

Aggregate Structure Identification

- G. Ramalingam et al

- Year 2000 problem
 - Used to identify date type variables
 - Made maintenance easier
- Improving static analysis algorithms
 - Aggregates considered like scalars
 - Imprecision creeps in
 - Do analysis in terms of atoms
 - Precision improves!

Aggregate Structure Identification - Example

A.

```
int F1,F2,F3,F4;
```

B.

```
int [4];
```

C.

```
int F5,F6,F7,F8;
```

RESULT.

```
int;
```

```
move 17 to F1;
```

```
move 18 to F2;
```

```
move A to B;
```

```
move B to C;
```

```
move F5 to RESULT;
```

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```

A

4

A light blue rectangular box containing the number 4, with a black arrow pointing from the text 'A' to its top-left corner.

B.

```
int [4];
```

B

4

A light blue rectangular box containing the number 4, with a black arrow pointing from the text 'B' to its top-left corner.

C.

```
int F5,F6,F7,F8;
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RESULT.

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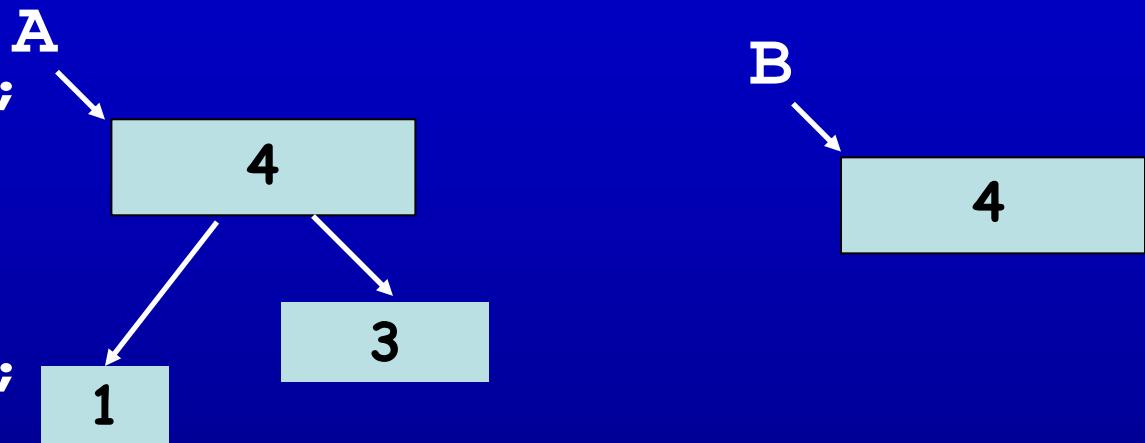
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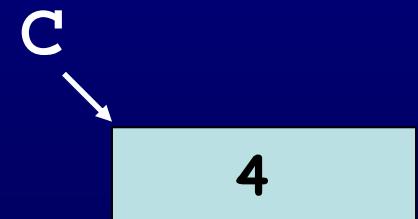
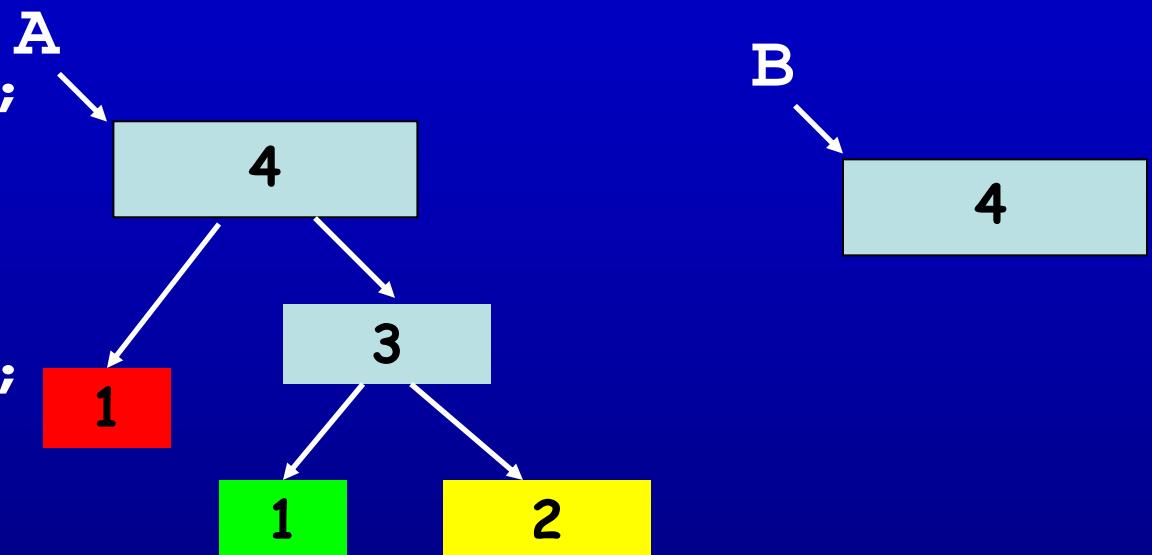
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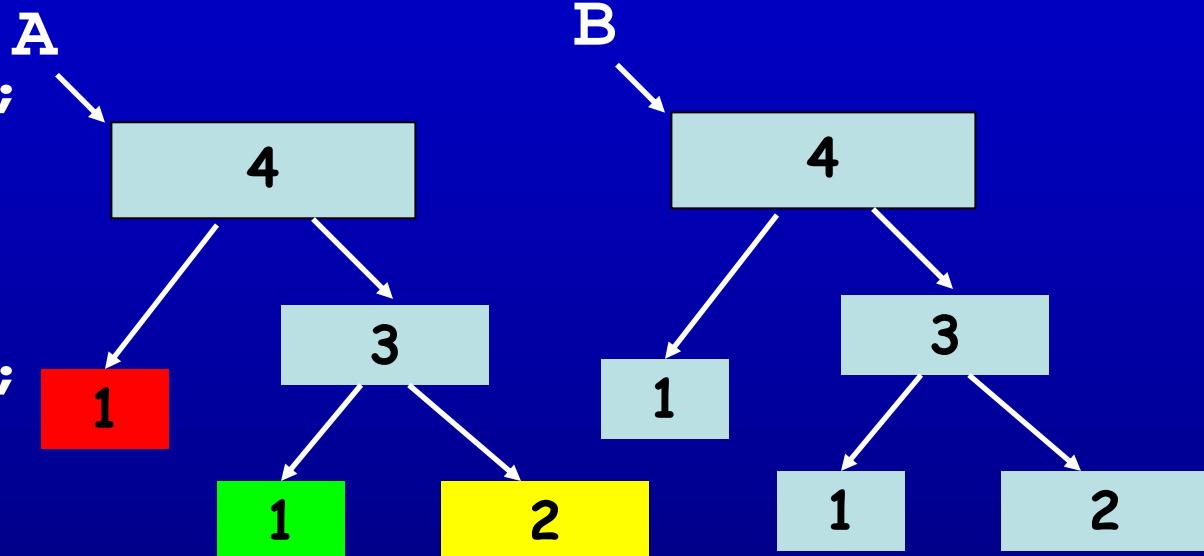
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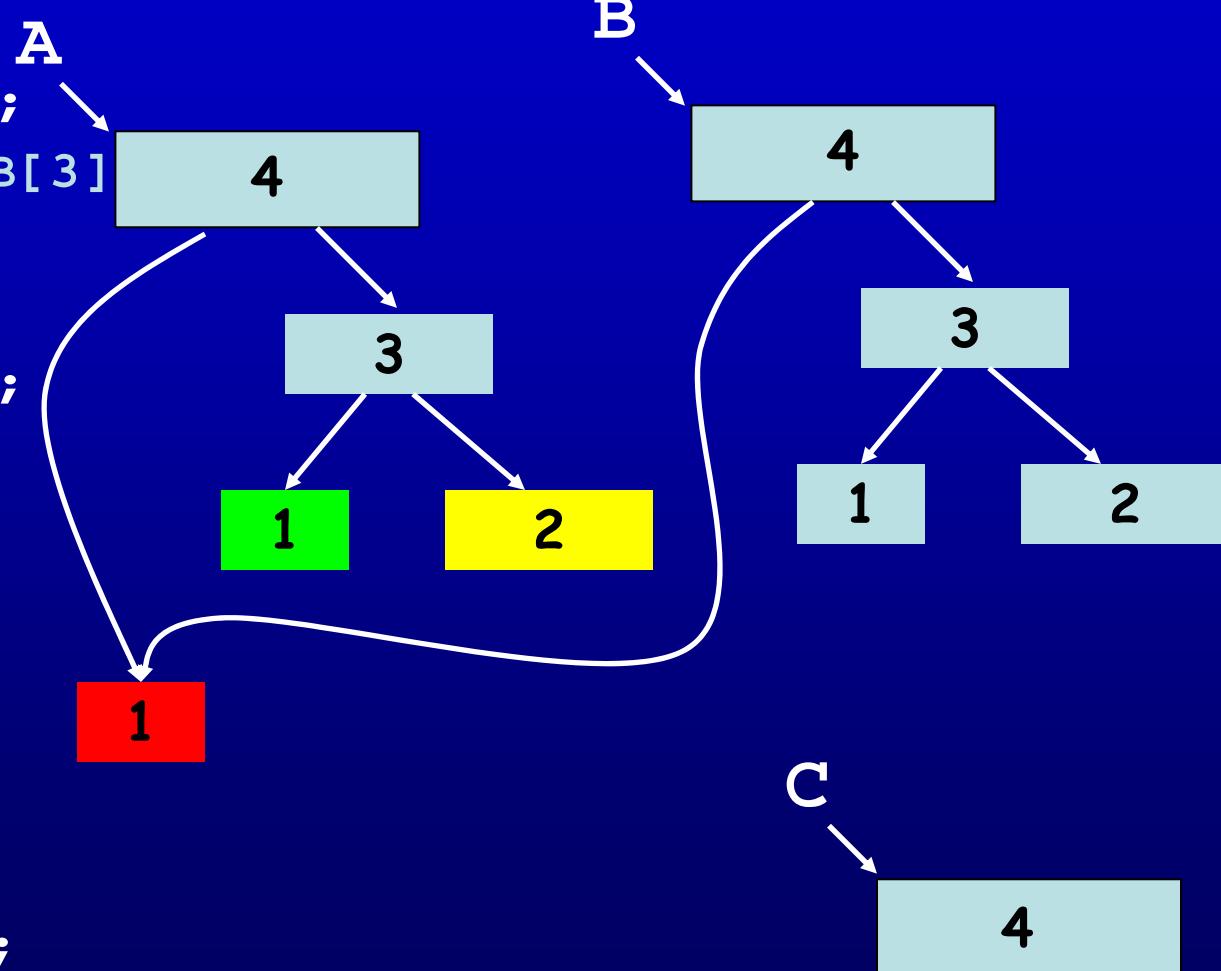
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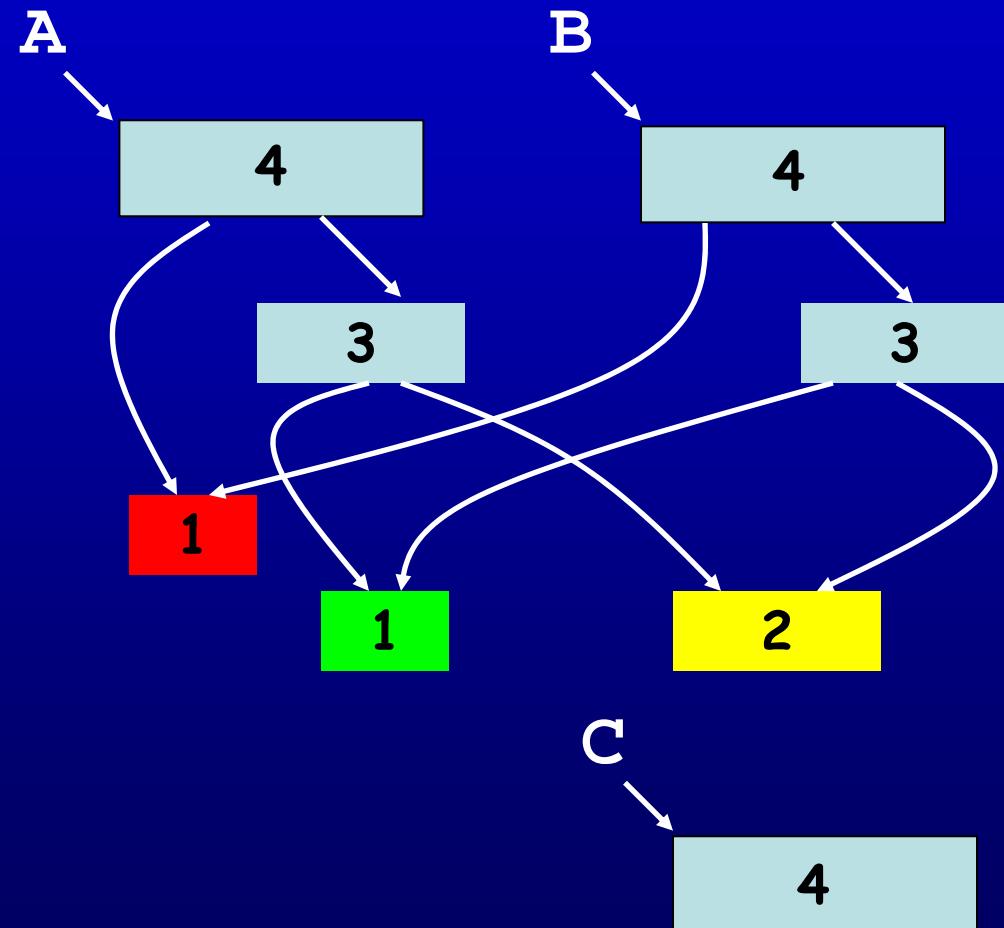
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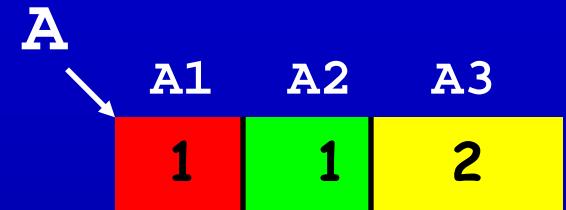
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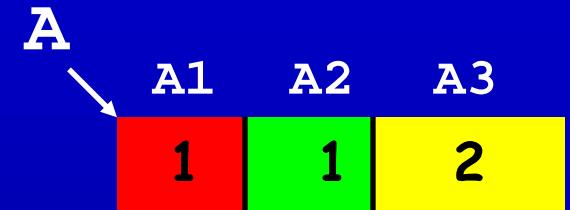
```
move 17 to F1; {17 -> A1}
```

```
move 18 to F2; {18 -> A2}
```

```
move A to B; {(A1,A2,A3)->(B1,B2,B3)}
```

```
move B to C; {(B1,B2,B3)->(C1,C2,C3)}
```

```
move F5 to RESULT; {C1 -> RESULT}
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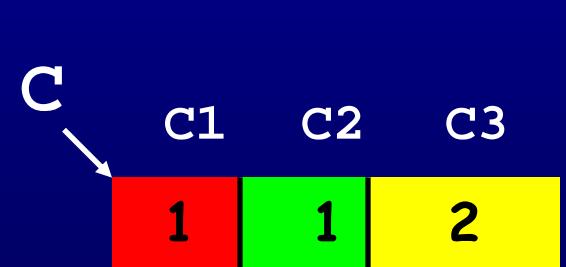
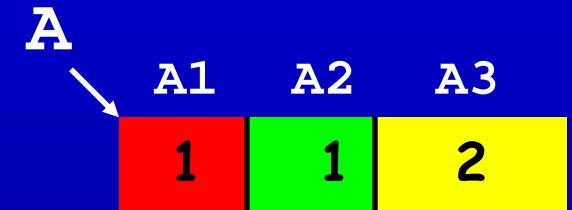
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```
move F5 to RESULT; {C1 -> RESULT}
```



Aggregate Structure Identification in Object Code

Activation
Records



- These areas of memory
 - Aggregates in the algorithm

Aggregate Structure Identification in Object Code

Activation
Records



- These areas of memory
 - Aggregates in the algorithm
- Identify the structure

Aggregate Structure Identification in Object Code

Activation
Records



- These areas of memory
 - Aggregates in the algorithm
- Identify the structure
- Use the atoms of the aggregates as the entities

Minilanguage

- Input to the atomization algorithm
- Getting the minilanguage program
 - Retain only data transfer instructions
 - DataRef = DataRef
- DataRef
 - Data reference - three kinds
 - Program Variables
 - Range - for fields of aggregates
 - mov 17 to F1
 - $A[1:4] = \text{int_const}[1:4]$
 - Statically indeterminate element of an array
 - mov 12 to B[i]
 - $B[1:16]\backslash 4 = \text{int_const}[1:4]$

Generating the Minilanguage file

- Which part of an aggregate is read/written?
- Clear in high level languages
- Not evident in object code

```
int main(){
    int a[10],i,j;
    j=0;
    for(i=0;i<10;++i){
        a[i]=i;
    }
    return j;
}
```

a[1:40]\10=i

```
; ebx corresponds to variable i
sub    esp, 44
mov    [esp+40],0 ; j = 0
xor    ebx, ebx ; i = 0
lea    ecx, [esp]
loc_9:
    mov    [ecx], ebx ; a[i]=i
    inc    ebx ; i++
    add    ecx, 4
    cmp    ebx, 10 ; i<10?
    jl    short loc_9 ;
```

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    cmp    ebx, 10 ; i<10?
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```

????=ebx[1:4]

Generating the Minilanguage file

- Which part of an aggregate is read/written?
- Inferred from the linear relationship among registers

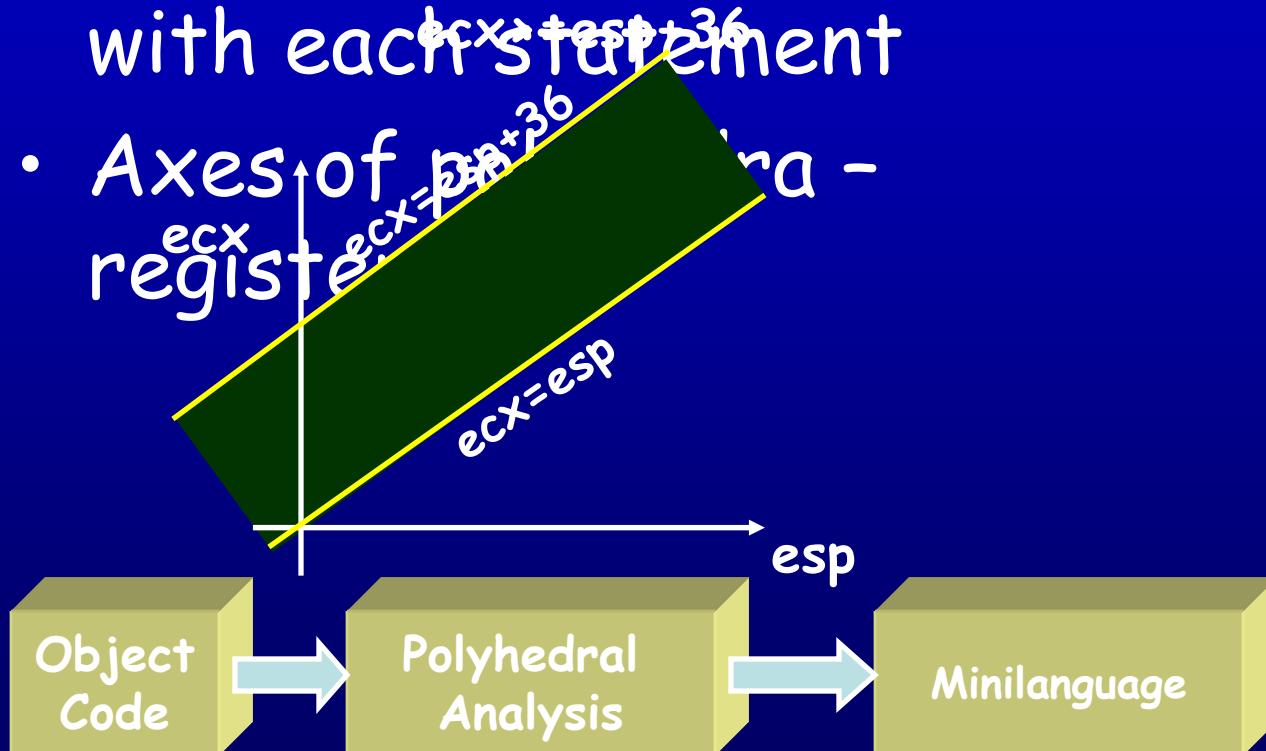
```
int main(){
    int a[10],i,j;
    j=0;
    for(i=0;i<10;++i){
        a[i]=i;
    }
```

$ecx >= esp \&&$
 $ecx <= esp + 36$
 $\therefore AR[1:40] \setminus 10 = ebx[1:4]$

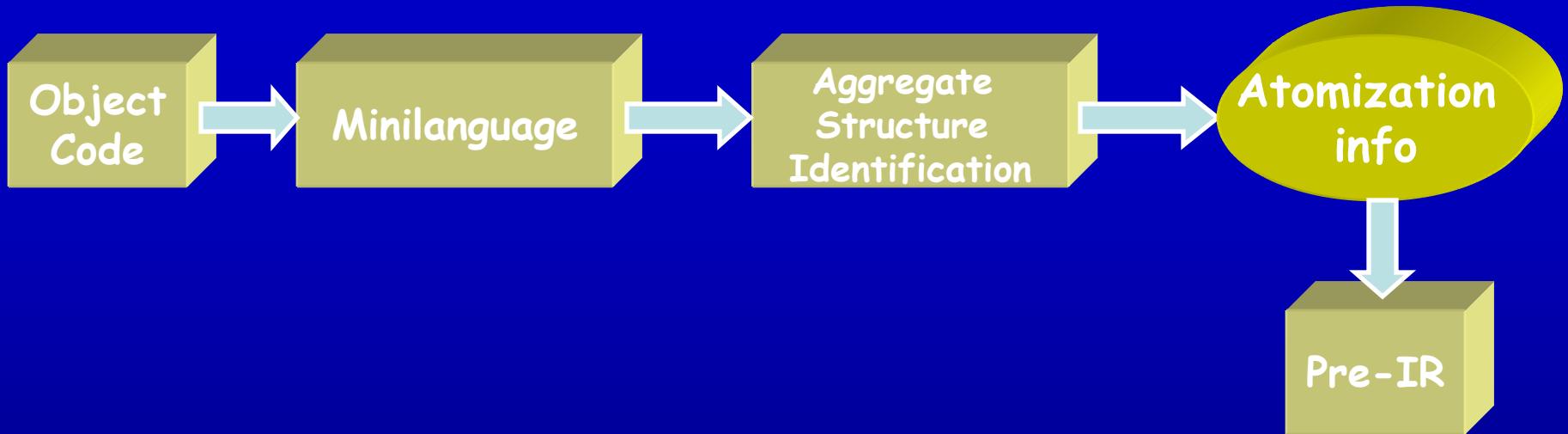
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        inc    ebx ; i++
        add    ecx, 4
        cmp    ebx, 10 ; i<10?
        jl    short loc_9 ;
```

Linear Relationships among Registers

- Use convex polyhedra
- Associate a polyhedra with each statement
- Axes of polyhedra -



Steps in the Algorithm



- Object Code → Minilanguage (Polyhedral analysis)
- Feed minilanguage to Ramalingam's analysis
- Identify the atoms
- Create pre-IR
- Feed it to codesurfer

Demo

```
struct Point{  
    int x,y;  
};  
struct Point g_pt={10,20};  
int gl_int=100;  
  
int main() {  
    struct Point l_a_pt[10];  
    int i;  
  
    g_pt.x=gl_int;  
    for(i=0;i<10;++i) {  
        l_a_pt[i].x=g_pt.x;  
        l_a_pt[i].y=g_pt.y;  
    }  
    return 0;  
}
```

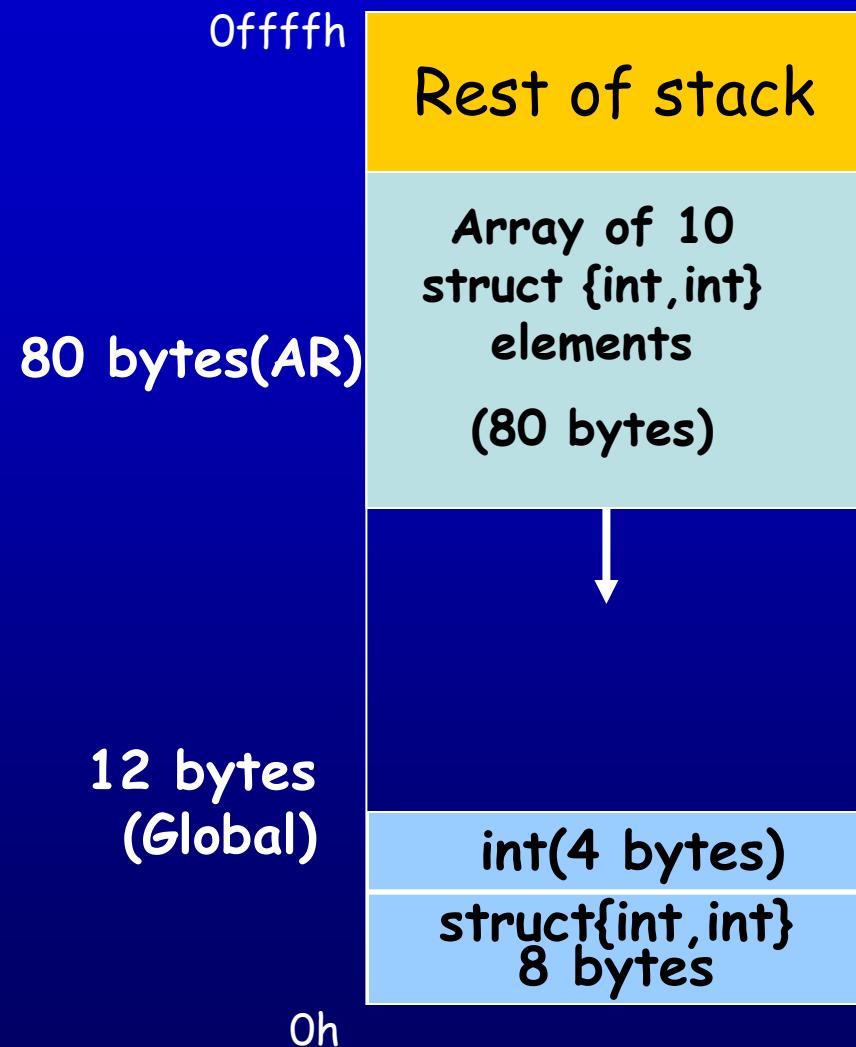
```
public _main  
    mov    edx, ds:gl_int  
    sub    esp, 50h  
    mov    ds:g_pt@x, edx  
    lea    eax, [esp+50h+var_4C]  
    mov    esi, ds:dword_4  
    push   esi  
    mov    ecx, 0Ah
```

```
loc_2B:  
    mov    [eax-4], edx  
    mov    [eax], esi  
    add    eax, 8  
    dec    ecx  
    jnz    short loc_2B  
    pop    esi  
    add    esp, 50h  
    retn
```

_main

endp

Demo



```
public _main
    mov    edx, ds:gl_int
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        add    eax, 8
        dec    ecx
        jnz    short loc_2B
        pop    esi
        add    esp, 50h
        retn

_main
endp
```

minilanguage file

```
decl eax 4;  
decl ecx 4;  
decl edx 4;  
decl ebx 4;  
decl esp 4;  
decl ebp 4;  
decl esi 4;  
decl edi 4;  
decl _main_AR 84;  
decl Global 12;  
decl const 4;
```

```
edx[1:4] = Global[9:12];
```

```
Global[1:4] = edx[1:4];
```

```
esi[1:4] = Global[5:8];
```

```
ecx[1:4] = const[1:4];
```

```
_main_AR[5:84]\10[1:4] = edx[1:4];
```

```
_main_AR[5:84]\10[5:8] = esi[1:4];
```

Demo

public _main	
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sub	esp, 50h
mov	ds:g_pt@x, edx
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mov	[eax-4], edx
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add	eax, 8
dec	ecx
jnz	short loc_2B
pop	esi
add	esp, 50h
retn	
	endp

_main

Conclusions

- No properly defined entities in object code
- Ramalingam's atomization algorithm
 - Atoms can be the entities
- Now, existing static analysis algorithms can be adopted to object code

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