



Reaching New Scales with the CMS HTCondor Global Pool



JAMES LETTS & ANTONIO PÉREZ-CALERO on behalf of the Submission Infrastructure Group of the CMS Experiment

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Why a Global Pool?

- Before the Global Pool, we had different HTCondor pools for data analysis and production activities.
- As we moved away from the original tiered LHC Computing Model, we wanted to flexibly run different types of workflows across tiers, as well as integrating new types of Cloud and allocation-based resources.

Tiered (MONARC) Model

when networking was a scarcer resource

Global Model

- We now have a unified, multi-core, HTCondor Global Pool with re-usable pilots.
- For reasons of stability for CMS data taking, the Tier-0 has its own separate pool.
- CMS is also moving away from data locality: AAA, xrootd, caching, etc.
- Makes resource scheduling much more complicated! CPU, Memory, Disk, and now I/O. Network is still scarce in places.
- In 2016, CMS moved into a resource-constrained environment for the first time.

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Auto-clusters

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Challenges

- Most recent scalability limits over the past few years have been found in the Central Manager.
- Scale tests with single-core jobs and conducted with the <u>OSG in 2014</u> at 200,000 static slots found that separation of the CCB's onto hardware separate from the Central Manager was essential to go beyond 150,000 CPUs (in the lab).

Challenges

- However, in the wild, a new blocker arose in 2016 at ~155,000 CPUs in a multi-core environment.
- Symptom: Central Manager machine dropping UDP updates.
- CMS worked closely with the HTCondor developers to study the problem. Finally found out that the Top Collector was being blocked by **queries**.
- Limited number of forked query workers exhausted, remaining queries were blocking the Top Collector.

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Solution

• Developers' solution was to queue updates, not let them go to the Top Collector, and prioritize queries from the Negotiator.

Evolution

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Earlier Issues

- Negotiator scalability has been an issue we have encountered several times so far.
- Since the beginning of 2015 we run 3 separate
 Negotiators in parallel based on resource type:
 Tier-1, U.S. Tier-2, and the rest. Each group
 ~80K CPU cores.
- Allows us to do resource-based fair share: production gets 95% of the Tier-1 sites, while rest are 50% physics analysis.
- HTCondor developers have been parallelizing the Negotiator even more since.

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Scale Tests

- In principle, CMS wants to find and fix blockers in scale tests, not in a production system.
- Last round of scale tests with OSG in 2014 used the concept of "<u>über-glideins</u>": one pilot launches multiple startd's, thus achieving the I/O of a much larger pool:

 Using a factor of 32, can reach the I/O of 500K startd's using only 15,625 physical cores.

2017 Scale Tests

- In principle CMS can do something similar within the existing Global Pool, since we are not worried about the scalability of glideinWMS so much as HTCondor.
- Planning the next round of tests for August, which is historically a low period of usage after the major summer conferences, on new, beefier VMs provided by CERN/IT (~96GB RAM).
- Hope to find the next blockers in a test environment rather than in the wild!

Future

- Future challenges besides global scalability that we are facing for 2017-2018:
- Scheduling I/O: Now that individual sites are approaching 50-100K CPU cores, how can we not kill the network or the storage at the site? Or even across groups of sites? (CMS is moving away from data locality)
- Improving scheduling efficiency (next slide):

Improving Scheduling Efficiency

- Filling multi-core p-slots is a multi-dimensional problem:
 - Job requirements (e.g. time, CPUs, memory, resizable jobs)
 - Bursty nature of job arrival (time)
 - Resource constraints (CPUs, Memory)
 - Fair-share and priority (ranking)
 - Pilot lifetime (time)

CPU	4							
1	job1		job5					
2			job6					
3	job2				job	8		
4	job3	ob3 job4		b7	j		b9	
								time

Improving Scheduling Efficiency

- Pool partitioning evolves over time to serve the demand.
- The challenge is that sometimes more CPU cores than desired are left unscheduled.

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Improving Scheduling Efficiency

- There is a ~5% irreducible amount of wasted CPU from retiring glideins (p-slots). Tunable?
- High-memory jobs can take all of the RAM of a pilot, so that (justifiably) some CPU will be left unused.

Conclusions

- We thank the HTCondor development team for their close collaboration.
- We have met some interesting scalability and stability challenges over the past couple of years and look forward to reaching even greater heights in the years to come.